

One Key To Rule Them All

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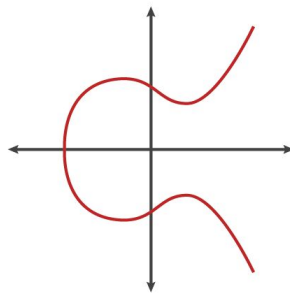
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Refresher: Elliptic curves

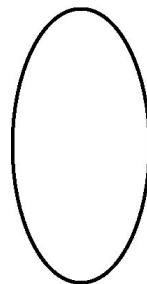
Seen last year at NSec !

I DON'T UNDERSTAND WHY
PEOPLE GET CONFUSED ...
I DON'T LOOK ANYTHING
LIKE YOU!



ELLIPTIC CURVE

I THINK IT'S THE NAME!
LET'S ASK JAVA AND
JAVASCRIPT TO SEE HOW
THEY DEAL WITH IT



ELLIPSE

In the talk “Getting ahead of the elliptic curve” by Martijn Grooten

An elliptic curve will simply be the set of coordinates (x,y) :

$$\{(x, y) \in \mathbb{F}_p^2 \mid y^2 = x^3 + ax + b, \quad 4a^3 + 27b^2 \neq 0\} \cup \{0\}$$

on which we can define an addition law “+” that is well-behaved.

Crypto refresher: ECC

ECC aka “Elliptic Curve Cryptography”

- Working with the subgroup of prime order m of an **elliptic curve C** generated by a given **generator G**
- Security based on the hardness of the EC discrete logarithm problem (DLP):

If the point $Q = nP$ is obtained by scalar multiplication of an integer n with the point P , being given only Q and P , it is hard to retrieve n .



Refresher: EC

“Scalar multiplication”?

Comes naturally from the additive operation “+” under which Elliptic Curves over a finite field \mathbb{F}_p form an abelian group, as you’ve seen last year ;)

$$nP = \underbrace{P + \dots + P}_{n \text{ times}}$$

By definition, scalar multiplication is naturally distributive over EC addition:

$$aP + bP = (a + b)P$$

Crypto refresher: ECC

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- Security based on the hardness of the EC discrete logarithm problem.

In common EC schemes (ECDSA, ECDH, ECIES), we usually have that:

- **A private key** is simply an integer $k < p$ (it is an element of \mathbb{F}_p)
- **A public key** is a point $P = (x, y)$ on the curve generated by **scalar multiplication** of the private key with the base point: $P = kG$
 - where (x, y) are the coordinates of the point on the curve C

Using ideas coming from the Bitcoin world

Have you ever heard of the cool idea behind the [BIP-32](#) proposal also known as “**Hierarchical Deterministic Wallets**” by Pieter Wuille?

It introduced the idea of having so-called *child keys* derived from *parent keys* when working with a EC signature algorithm such as ECDSA or Schnorr signatures.

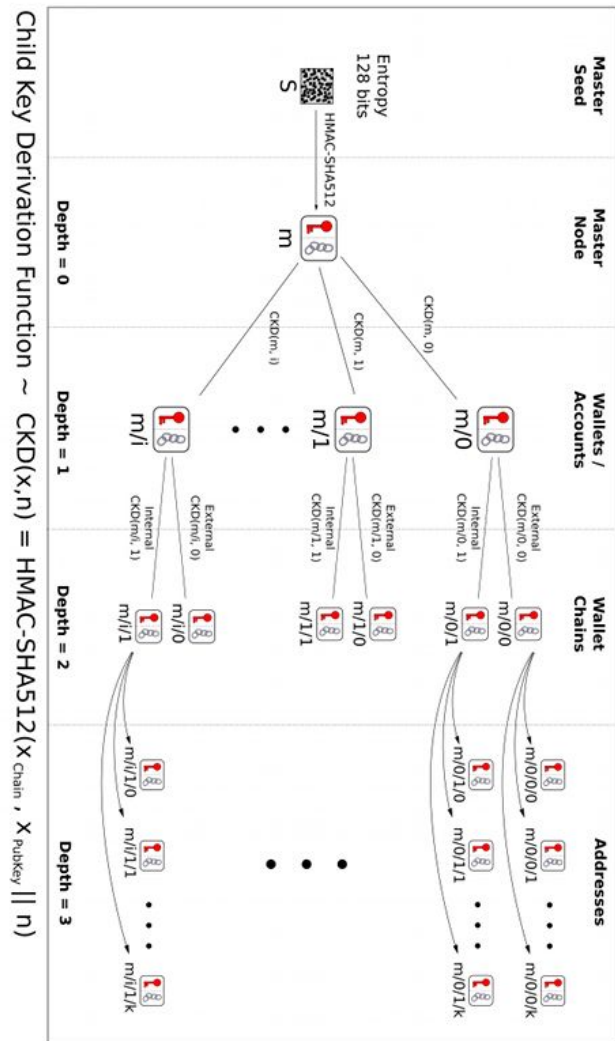
This is made possible using the simple math on EC we just saw!

The principle behind BIP-32

In Bitcoin, addresses are **the hash of a ECDSA public key**, which is a point on an Elliptic Curve. It uses the same kind of key pair that we discussed earlier.

The idea is that if you have a private key k and a public key $P = kG$, you can pick a deterministically generated value r and if you take the point $Q = P + rG$, you actually just computed the public key for the private key $k + r$, since

$$P + rG = kG + rG = (k + r)G \text{ by distributivity !}$$



BIP-32, the stuff I skipped

- It uses “extended” keys that contain both k and a so-called “chain code”
- It uses HMAC to derive the value r using the chain code as a HMAC key
- It uses indices i to index its child keys
- It has a notion of hardened keys that do not allow for public key derivation
- Its keys are nodes of a key tree because it enables fun bitcoin stuff

Nonetheless, the core idea is that $\mathbf{P} + r\mathbf{G} = (\mathbf{k} + r)\mathbf{G}$ and that’s reusable!

Public keys... for what?

- Connecting to remote machines without relying on ~~weak~~ mere passwords!
- Sending PGP emails?
- Connecting to ~~cool~~ modern VPNs such as **Wireguard**.
- Sending / receiving cryptocurrencies.
- Encrypting passwords with tools such as **gopass**.
- ...



ECC SSH keys out there

Most common elliptic curves for SSH according to a world-wide key collection:

- **secp256r1** 97,68% (aka “NIST P-256”)
- secp521r1 1,87%
- Curve25519 0,37% (Ed25519)
- secp384r1 0,07%



Ed25519 isn't directly compatible with BIP-32 because it hashes its private key. (Can be worked around by tweaking the signing algorithm, but not a drop-in replacement.)

What for?

SSH key derivation sounds fun, but why is it interesting?

- Anyone can generate a new public key for you being given your public key, and you'll have the private key as long as they tell you the random r they used
- You can generate new keys as you wish, but just need to store one single key!
- It's ~~easy~~ easy to use, and doesn't weaken security
- It might allow for fun PKI systems, where you can link a public key to another one by revealing the random value r shall

Demo

`./1key -key path/to/key derivation-code`

-a allows to export your private key (and even to encrypt it)

-q adds the key to the ssh-agent on unix devices and then quit

-rm removes ALL the keys in ssh-agent. But `~/.ssh/*` keys will be reloaded.

-s "user@192.168.0.42" will "exec" the ssh command with the provided args.

Disclaimer: just a POC

But you can test it using your keys, and it works!

I have open-sourced my (go) code on:

github.com/kudelskisecurity/1key

How about security?

Having a single SSH key is not bad *per se*. However it can leak data:

- <https://github.com/anomalroil.keys>
- <https://gitlab.com/anomalroil.keys>

Using child keys is not worse than using a single private key.

It can also enable some nice “recovery” possibilities, if you loose your computer.

BTW: always store your private keys in an encrypted form!

But...

From child key to master key!

A child key is simply the master key k plus a given value r .

So, if one knows the child key $(k + r)$, and knows the value r , it is easy to recover the value $k = (k + r) - r$

This can both:

- allow you to have different private keys on each of your devices and yet a single “master public key”
- be a problem if one of your private child key is compromised and the attacker knows the value r

(Can be done using `./1key -i -r “hex-value”`)

Applying this idea to other fields?

- Wireguard key distribution? Uses Ed25519, where the private key is hashed...
Requires some tweaking on Wireguard itself :(
- Honest ransomware aka “RIP-001” ;)
- “Disposable keys” with “opt-in” non-repudiability (you can reveal r or not)

Conclusion

You can try it out, if you want. Open issues on Github, or submit PR!

Mind your keys, keep them encrypted!

Thank you!



Follow me on Twitter:
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Links

- Just a reminder to check your RSA public keys:
<https://keylookup.kudelskisecurity.com/>
- Find my open source code on Github:
<https://github.com/kudelskisecurity/1key>
- Find more results and analysis on our blog:
<https://research.kudelskisecurity.com>
- Download these slides on my homepage (in a couple of days):
<https://romailer.ch/page/talks/>